

Kerberos

Secure and efficient authentication and key distribution

Johannes Löttsch and Meike Zehlike

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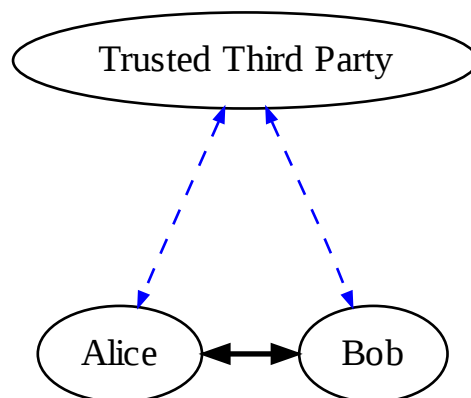
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Use cases for Kerberos / Alternatives

Trusted Instance for:

- ▶ Keydistribution
- ▶ Authentication
- ▶ SingleSignOn



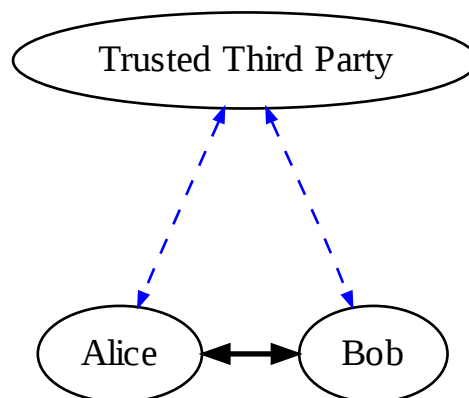
Well known alternative Trusted Third Parties:

- ▶ CA in hierarchical PKI
 - ▶ Asymmetric Cryptography \implies slow / expensive
- ▶ Members of Web of Trust (e.g. PGP)
 - ▶ Hard to say how trustworthy a „Trustpath“ is

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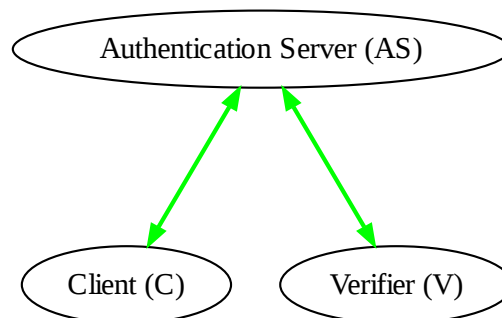


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Basic Protocol

- ▶ Based on symmetric Needham-Schroeder-Protocol
- ▶ Assumption: Each participant exchanged Key with AS on save channel $(K_{AS,C}, K_{AS,V})$

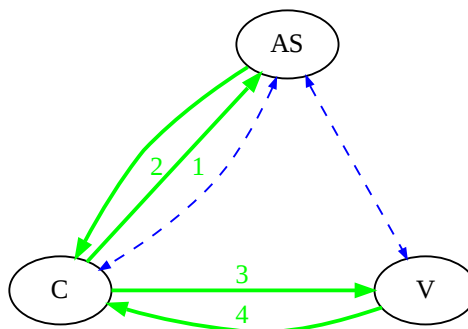


Procedure:

1. Authentication Request: $(C, V, time_{exp}, n)$
2. Authentication Reply: $E_{K_{AS,C}}(V, time_{exp}, n, K_{C,V}), E_{K_{AS,V}}(C, time_{exp}, K_{C,V})$
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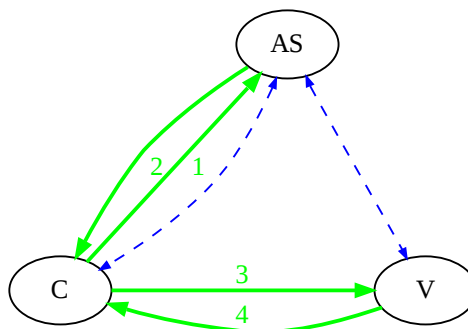


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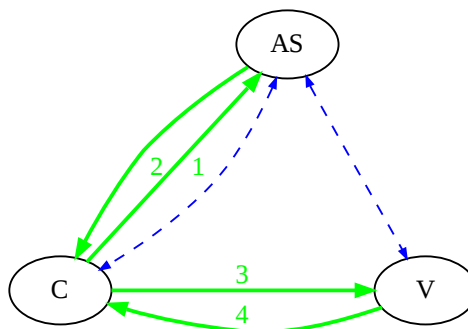


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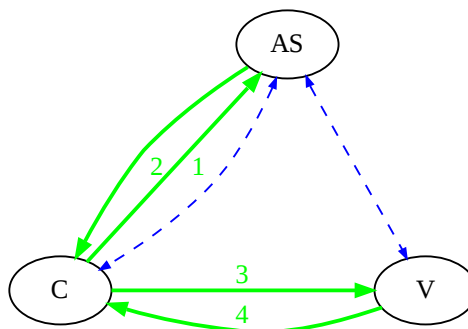


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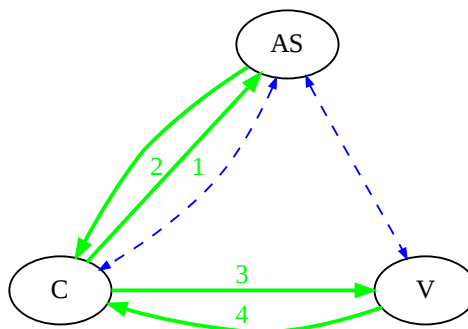


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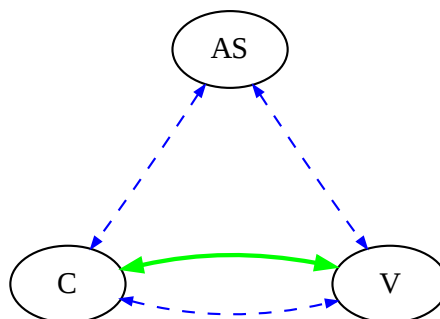
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Ticket Granting Tickets

- ▶ $K_{AS,C}$ needs to be handled with caution
- ▶ usage from different machines
 - ▶ should not be stored on clients \implies always created from a password

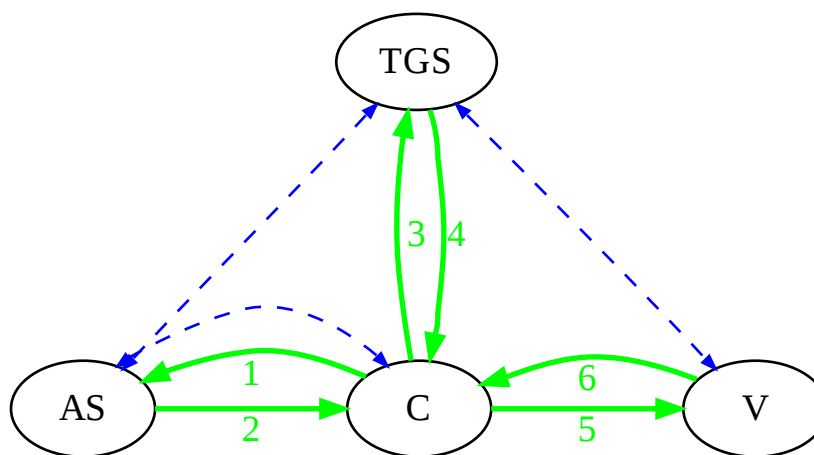
Howto cache credentials?

- ▶ 1-2: Authentication with AS using password \implies TGT
- ▶ 3-4: Authentication with TGS using TGT \implies SessionTicket for V
- ▶ 5-6: Authentication with V using SessionTicket
- ▶ 3-6 can be repeated until $time_{exp}$ of TGT

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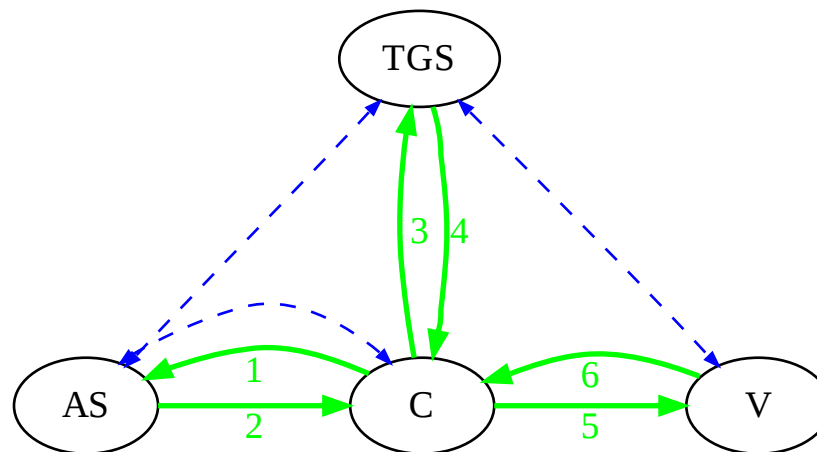


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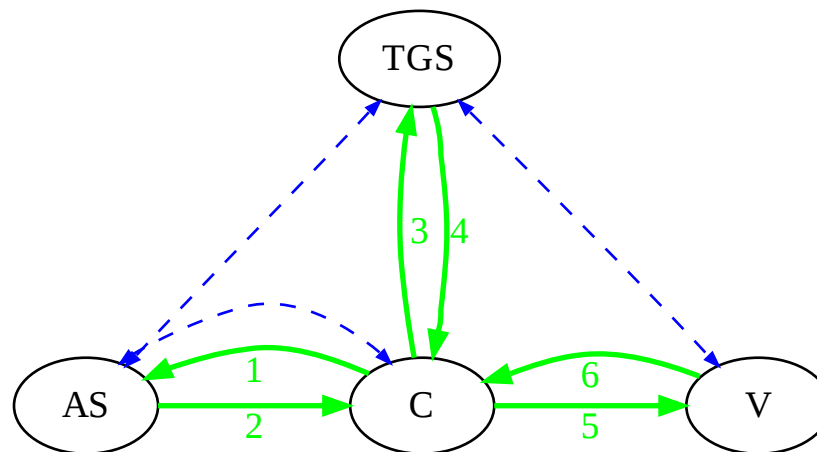


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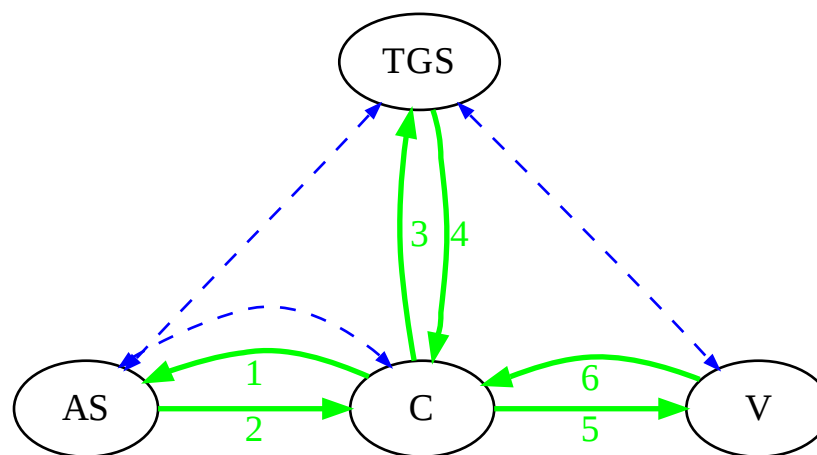


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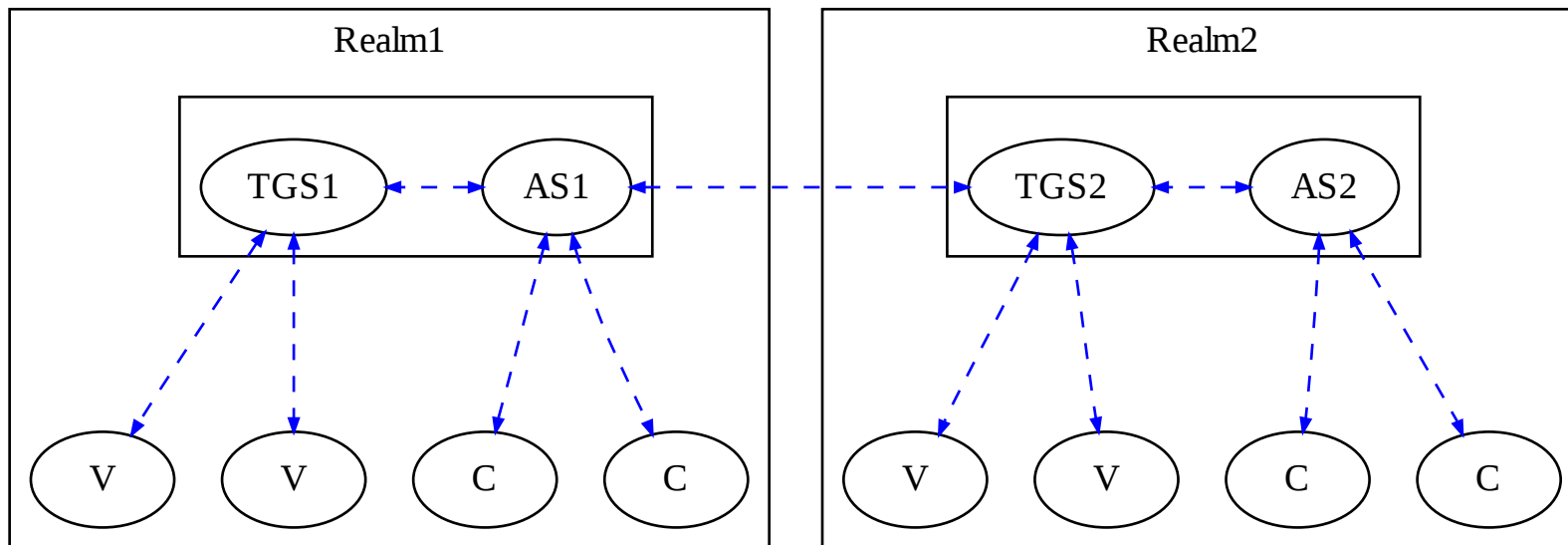
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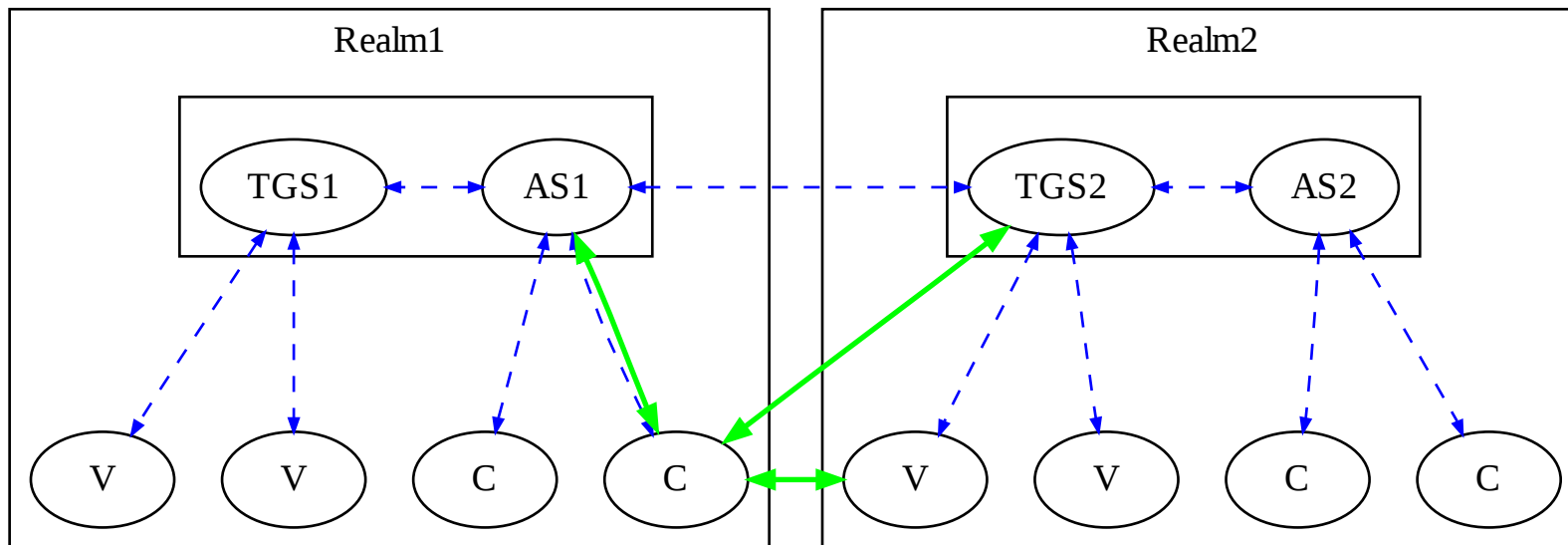


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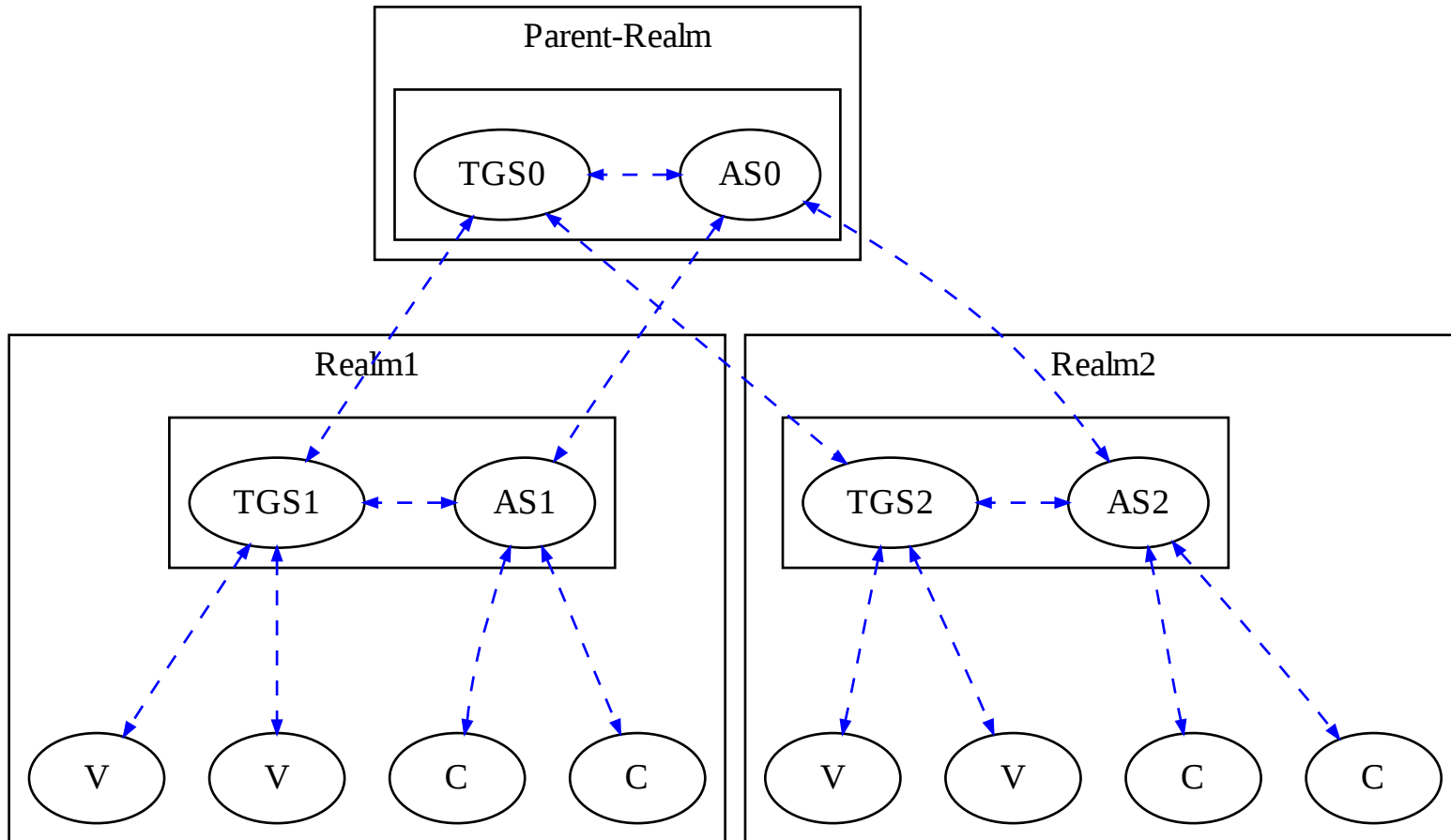
Cross-Realm-Authentication



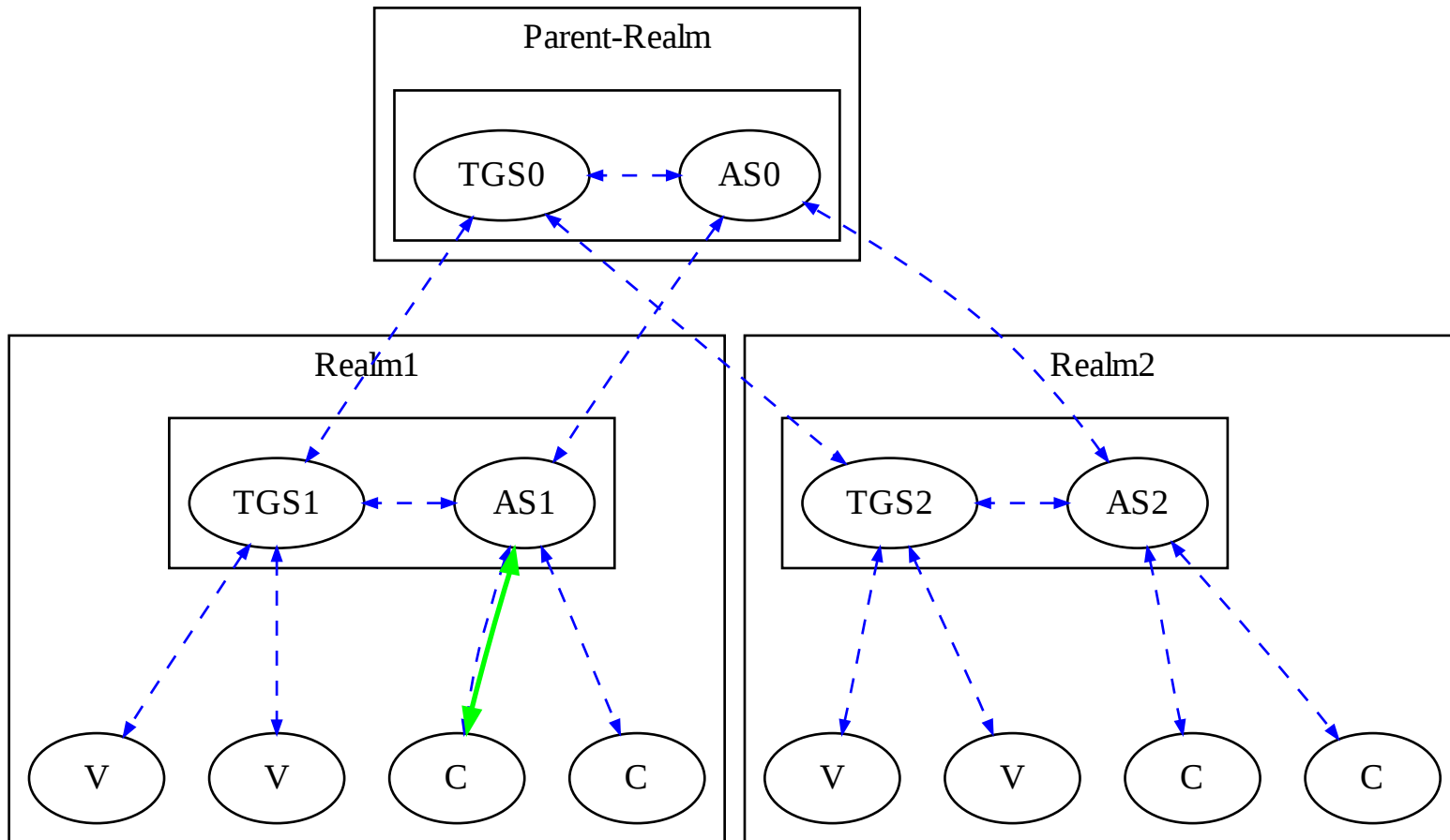
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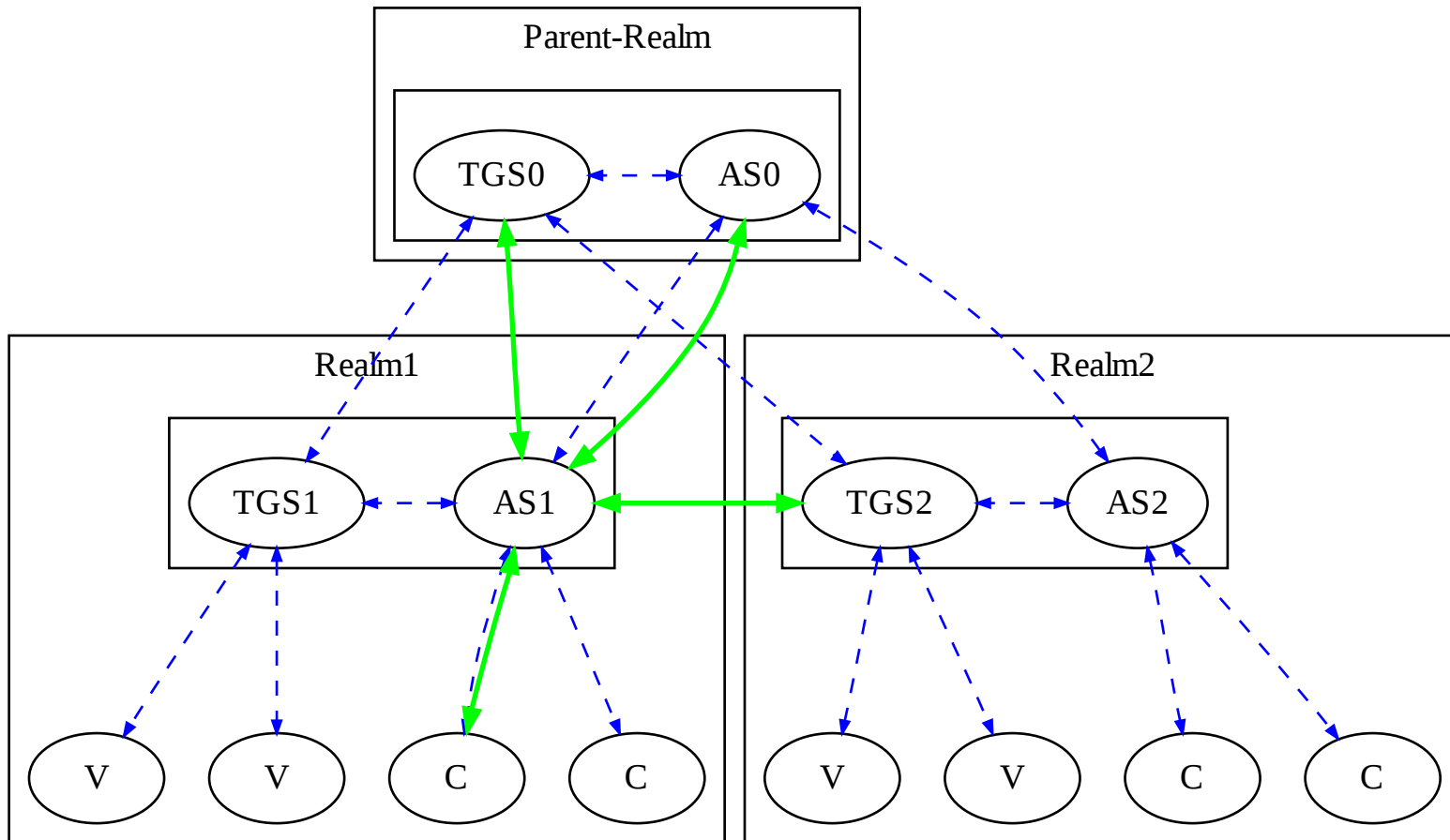
Authentication in hierarchic Realms



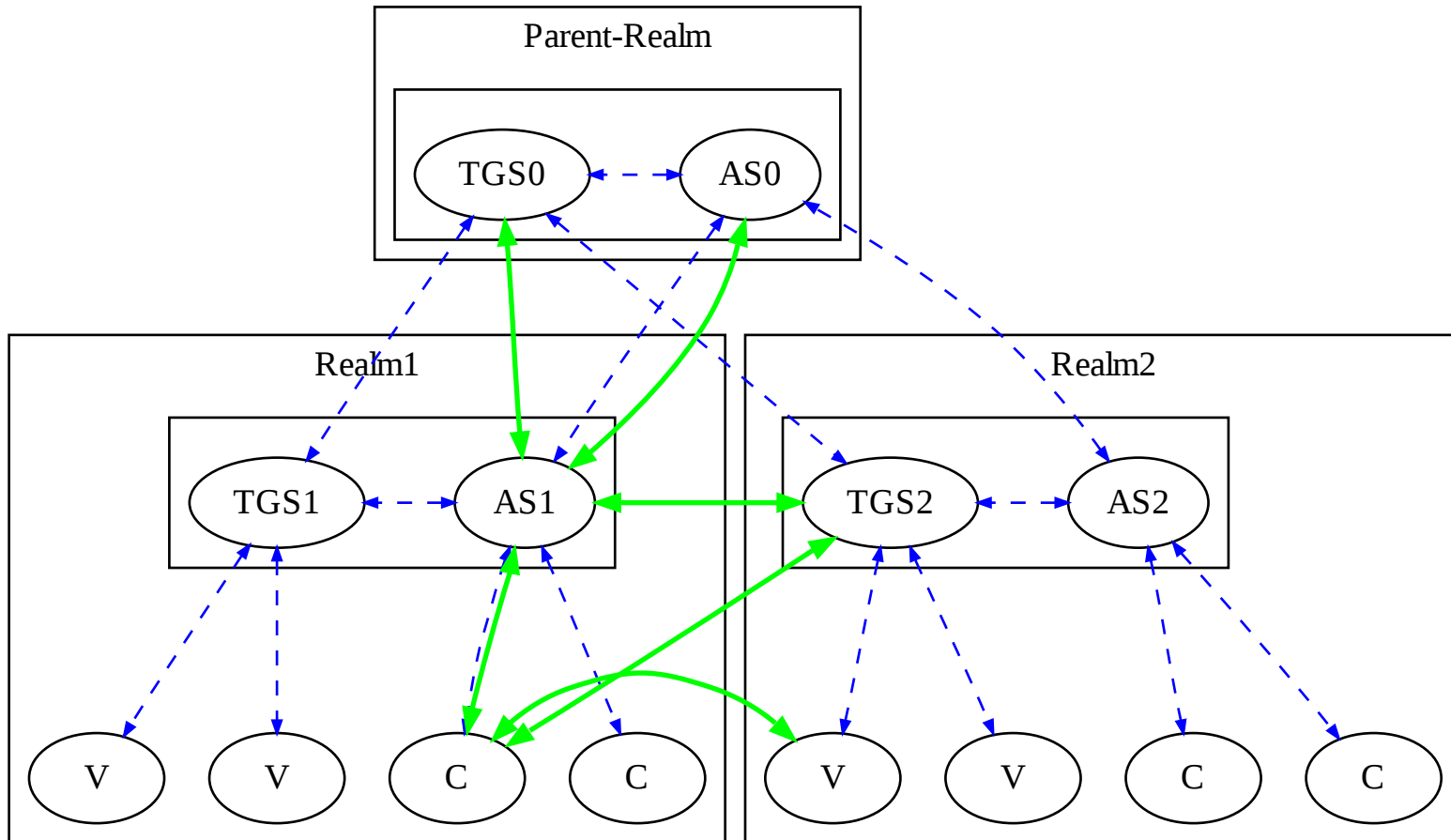
Authentication in hierarchic Realms



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Authentication in hierarchic Realms



Installation on Debian / Ubuntu

```
$ wget http://johannesloetzsch.de/about/kerberos/setup_krb.sh
```

```
...
```

```
read, UNDERSTAND, modify
```

```
...
```

```
$ chmod +x setup_krb.sh
```

```
$ sudo ./setup_krb.sh
```

Add new user (create $K_{AS,C}$)

```
$ sudo kadmin.local
```

```
kadmin.local: addprinc newuser
```

```
Enter password for principal "newuser@EXAMPLE.COM": $password
```

```
Re-enter password for principal "newuser@EXAMPLE.COM": $password
```

```
Principal "newuser@EXAMPLE.COM" created.
```

```
kadmin.local: list_principals
```

```
...
```

```
newuser@EXAMPLE.COM
```

```
...
```


kadmin (add_principal, ktadd, list_principals)

Add new server (create $K_{TGS,V}$)

kadmin.local: **addprinc -randkey host/servername.example.com**
Principal "host/servername.example.com@EXAMPLE.COM" created.

kadmin.local: **ktadd -k /tmp/key host/servername.example.com**
Entry for principal host/servername.example.com with kvno 3, encryption type
AES-256 CTS mode with 96-bit SHA-1 HMAC added to keytab
WRFILE:/tmp/key.

\$ sudo scp /tmp/key root@servername.example.com:/etc/krb5.keytab



Obtaining the TGT

```
$ kinit newuser
```

```
Password for newuser@EXAMPLE.COM: $password
```

```
$ klist -5
```

```
Ticket cache: FILE:/tmp/krb5cc_1000
```

```
Default principal: newuser@EXAMPLE.COM
```

```
Valid starting Expires Service principal
```

```
10/07/09 20:16:04 10/08/09 06:16:04 krbtgt/EXAMPLE.COM@EXAMPLE.COM
```

```
renew until 10/08/09 20:15:58
```

```
$ kdestroy
```

```
$ klist -5
```

```
klist: No credentials cache found (ticket cache FILE:/tmp/krb5cc_1000)
```

Use the Kerberos-principal for SSH

```
$ yes " |sudo adduser --disabled-password newuser
```

```
$ kdestroy
```

```
$ ssh newuser@servername.example.com
```

```
newuser@servername.example.com's password: password is disabled
```

```
Permission denied, please try again.
```

```
$ kinit newuser
```

```
$ Password for newuser@EXAMPLE.COM: $password
```

```
$ ssh newuser@servername.example.com
```

```
newuser@servername.example.com: $
```

Done — SingleSignOn works

```
$ klist -5
```

```
Ticket cache: FILE:/tmp/krb5cc_1000
```

```
Default principal: newuser@EXAMPLE.COM
```

```
Valid starting Expires Service principal
```

```
10/07/09 20:27:09 10/08/09 06:27:09 krbtgt/EXAMPLE.COM@EXAMPLE.COM
```

```
renew until 10/08/09 20:27:03
```

```
10/07/09 20:27:14 10/08/09 06:27:09
```

```
host/servername.example.com@EXAMPLE.COM
```

```
renew until 10/08/09 20:27:03
```

```
$ ssh newuser@servername.example.com
```

```
newuser@servername.example.com: $ exit
```

```
$ ssh newuser@servername.example.com
```

```
newuser@servername.example.com: $
```

```
$ kdestroy
```

Further literature

DIY

-  [http://en.wikipedia.org/wiki/Kerberos_\(protocol\)](http://en.wikipedia.org/wiki/Kerberos_(protocol))
-  <http://gost.isi.edu/publications/kerberos-neuman-tso.html>
-  <http://techpubs.spinlocksolutions.com/dklar/kerberos.html>
-  <http://linux.die.net/man/1/kerberos>
-  <http://tools.ietf.org/html/rfc4120>